A Chaos-based Direct-Sequence/Spread-Spectrum Communication Scheme

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Abstract—This paper proposes a novel direct sequence/spread spectrum (DS/SS) communication scheme, where bit duration is varied according to a chaotic behavior but always equal to a multiple of the fixed chip duration in the communication process. To recover exactly the data bits, the receiver needs an identical regeneration of both the pseudo-noise (PN) sequence and chaotic behavior.

I. PROPOSED COMMUNICATION SCHEME

The proposed chaos-based DS/SS communication [1] scheme is presented in Fig. 1(a). Basically, the scheme is built around variable-position pulse and PN sequence (VPP-PNS) generator as in Fig. 1(b).

A. VPP-PNS Generator

The pulse train at the output 1 is expressed as follows:

$$p(t) = \sum_{n=1}^{\infty} P_{T_c} (t - t_n),$$
 (1)

where the n^{th} pulse is generated at the instance

$$t_n = t_{n-1} + N_{n-1}T_c = t_{n-1} + \left| \frac{X_n}{m} \right| T_c,$$
 (2)

and its position is determined by the interval

$$t_n - t_0 = \sum_{k=1}^n \left\lfloor \frac{X_k}{m} \right\rfloor T_c = \sum_{k=1}^n \left\lfloor \frac{F^{(k)}(X_0)}{m} \right\rfloor T_c;$$
 (3)

 $P_{T_c}(t)$ is the rectangular pulse shaping function given by

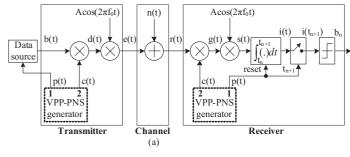
$$P_{T_c}(t) = \begin{cases} 1, & 0 \le t \le T_c, \\ 0, & otherwise. \end{cases}$$
 (4)

With the initial instance t_0 being fixed, Eq. 3 points out that this variation of position depends on the chaotic map F(.), initial value X_0 and count-step m.

B. Transmitter

In the transmitter, each pulse of the train p(t) triggers the data source to shift the next bit to the output. It means that the n^{th} bit is shifted at the instance t_n and its bit duration is equal to the interval of $[t_n,t_{n+1}]$. Based on Eq. 2, the duration of the n^{th} bit is determined by

$$T_{bn} = t_{n+1} - t_n = N_n T_c = \left[\frac{X_{n+1}}{m} \right] T_c = \left[\frac{F^{(n+1)}(X_0)}{m} \right] T_c.$$
 (5)



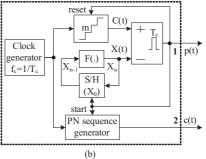


Fig. 1. (a) Proposed chaos-based DS/SS communication scheme, (b) Scheme for the VPP-PNS generator.

With a specific chaotic map F(.), in the iteration process, its output values vary chaotically in a defined range of $[X_{min}, X_{max}]$. This leads to the variation of the bit duration within a defined time range of $[T_{bmin}, T_{bmax}]$ in the communication process. Since the chip duration T_c is fixed, the number of chips per each bit, $N_n = T_{bn}/T_c$, called the spreading factor, also varies in a defined range of $[N_{min}, N_{max}]$ which is determined based on Eq. 5 as follows:

$$\begin{cases}
N_{min} = T_{bmin}/T_c = \lfloor X_{min}/m \rfloor, \\
N_{max} = T_{bmax}/T_c = \lfloor X_{max}/m \rfloor.
\end{cases} (6)$$

C. Receiver

The output signal i(t) of the integrator is reset to zero by the trigger of each pulse of the train p(t). Before each reset instance, the signal i(t) is sampled. The output value of the sampler at the instance t_{n+1} is determined by

$$i(t_{n+1}) = \int_{t_n}^{t_{n+1}} s(t)dt = \int_0^{T_{bn}} s(t)dt.$$
 (7)

Finally, the binary value of the data bit is recovered by the decision device based on the sign of the resulting sample.

D. BER Performance

Bit error rate (BER) performance for the proposed communication scheme with the channel noise being additive white Gaussian noise (AWGN) is evaluated approximately as follows:

$$BER = \sum_{N=N_{min}}^{N_{max}} P_N BER_N, \tag{8}$$

where P_N is the probability to the spreading factor being equal to N, and BER_N is the BER of the scheme for case of the spreading factor being equal to N. Assume that the chaotic values distribute uniformly in $[X_{min}, X_{max}]$. It means that the probability P_N is the same for all values of $N \in [N_{min}, N_{max}]$, so we have

$$P_N = 1/(N_{max} - N_{min} + 1). (9)$$

Based on the evaluation result of the error probability for the conventional DS/SS communication scheme with AWGN channel [2], the BER_N is determined by

$$BER_N = Q\left[\sqrt{2N\left(E_c/N_0\right)}\right],\tag{10}$$

here E_c is called the energy per chip, N_0 is the noise power spectral density, E_c/N_0 is known as the signal-to-noise ratio (SNR) per chip, and the function $Q[x] = \frac{1}{\sqrt{2\pi}} \int_x^\infty e^{-\frac{u^2}{2}} du$. From Eqs. 8, 9, 10, the theoretical evaluation of the BER performance is given by the following equation:

$$BER = \frac{1}{(N_{max} - N_{min} + 1)} \sum_{N=N_{min}}^{N_{max}} Q\left[\sqrt{2N\left(\frac{E_c}{N_0}\right)}\right]. \tag{11}$$

II. SIMULATION RESULTS AND DISCUSSIONS

A numerical simulation for the proposed communication scheme with AWGN channel are carried out in Simulink for cases of $[N_{min},N_{max}]=[26,36],\ [21,41],\ [53,73],\ [48,78].$ Here, the nonlinear function used is tent map [3].

BER performances obtained from the simulation as well as the evaluation according to Eq. 11 are presented in Fig. 2. It can be seen that the simulation and evaluation results are nearly the same. The BER performances for cases having the same average spreading factor (i.e., $N_{av} = (N_{min} + N_{max})/2$) are compared to each other and to that of the equivalent conventional scheme [2] which has a fixed spreading factor $(N = N_{av})$. We can find that the proposed scheme with [26, 36] case performs slightly worse than the conventional scheme with (N = 31) and better than the [21, 41] case. Similarly, the performance of [53,73] case is between that of the conventional scheme with (N = 63) and [48, 78] case. These comparison results point out a general rule is that the performance of the proposed system with $[N_{min}, N_{max}]$ gets better and approaches that of the corresponding conventional system with $(N=N_{av})$ when both the values N_{min} and N_{max} tend closer to the value N_{av} and vice versa.

Due to the sensitive dependence of the chaotic behavior on the initial value, a very slight error in its detection leads to exponentially increasing errors in the regeneration [3]. In this case, the detection of the variation of bit duration is completely incorrect, as a result, the BER is very high. This is proved by simulation results in Fig. 2. With only a very slight mismatch in the initial values (i.e., $\Delta X_0 = 0.0001$) between the transmitter and receiver, the BER performances for all four cases become the same and much worse than those of the no mismatch scheme. Although the ratio E_c/N_0 increases gradually, the mismatch BERs are nearly unvaried and approximately equal to $1.6\ 10^{-1}$. It is noted that the set of values of m, X_0 and F(.) is considered as a secret key. It is impossible for the intruder to recover correctly the data without having full information on this secret key.

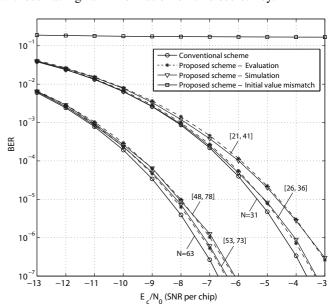


Fig. 2. Evaluation and simulation BER performances of the proposed communication scheme for different cases of $[N_{min}, N_{max}]$ in comparison to those of the equivalent conventional scheme.

III. CONCLUSION

Analysis and simulation results point out that the proposed scheme not only still maintains an approximately good performance but also achieves a significant improvement of the information security compared to the equivalent conventional scheme.

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